

# Compilation and Program Analysis

*Based on the lectures of  
Gabriel RADANNE and Ludovic HENRIO  
Notes written by Hugo SALOU*



*September 8, 2025*

# Contents

<b>1</b>	<b>Introduction.</b>	<b>4</b>
1.1	The RISC-V architecture. . . . .	5
1.2	Lexical Analysis. . . . .	5

Before we start, some administrative informations: the class is on Monday at 10:15AM and the tutorials are on Tuesday at 10:15AM. Tutorials will sometimes be about the notions from the day before.

# 1 Introduction.

A *compiler* is a program that'll given some code, return either an error or some other code in another language. However, compilation is not *only* about code generation: a large number of compilation techniques are not linked to assembly production. Moreover, languages can be:

- ▷ interpreted (*e.g.* Python sometimes);
- ▷ compiled into an intermediate language that will be interpreted (*e.g.* Java);
- ▷ compiled into another high level language (*e.g.* OCaml can be compiled into JavaScript);
- ▷ compiled “on the fly” (*e.g.* Julia or Python sometimes);
- ▷ several of the above.

A *compiler* will translate a program  $P$  into a program  $Q$  such that, for all entry, the output of  $Q$  is the same as the output of  $P$ . An *interpreter* is a program that, given a program  $P$  and an entry  $x$ , computes the output of  $P$  on  $x$ . This can be seen, in a way, we swap two quantifiers:

$$\underbrace{\forall P, \exists Q, \forall x, \dots}_{\text{compiler}} \quad \text{and} \quad \underbrace{\forall P, \forall x, \exists s, \dots}_{\text{interpreter}}$$

The quality of a compiler can be measured on multiple factors: its correctness, the efficiency of the generated code, its own efficiency. We will also touch on program analysis.

The goal of the labs will be to write a compiler for the RISC-V architecture. This part will be done in Python.

## 1.1 The RISC-V architecture.

RISC-V is an open-source architecture that is extensible. One of the basic components of RISC-V are *registers*. We can manipulate registers with operations such as `add`<sup>1</sup> or `addi`.<sup>2</sup>

We can also do *branching* in RISC-V in two kinds:

- ▷ *unconditional branching* ( $\rightsquigarrow$  jump and link) is done with `jal`;
- ▷ *test and branch* ( $\rightsquigarrow$  branch if lower than) is one with `blt`.

The first one is used to implement functions, and the other one (and variants) is used to implement an `if`.

All the details of the RISC-V operations can be found at:

[https://github.com/Drup/cap-lab25/blob/main/course/riscv\\_isa.pdf](https://github.com/Drup/cap-lab25/blob/main/course/riscv_isa.pdf).

We have an assembly language. This will be the last part of our compiler.

## 1.2 Lexical Analysis.

Lexical Analysis breaks down the code in tokens, known as *lexems*. Here, we use *regular expressions*. In our case, the tool we will use for our compiler is ANTLR.

---

<sup>1</sup>Adds the data from two registers into another register.

<sup>2</sup>Adds the data from one register with a constant, into another register